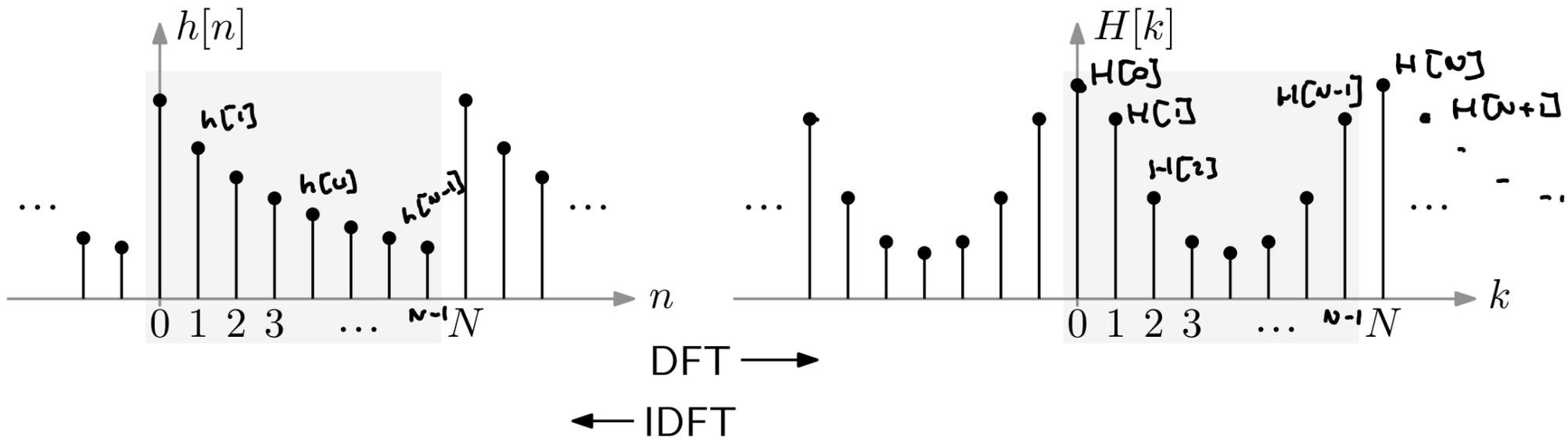


Fast Fourier transform (FFT)

And examples of how to compute things quickly

Herman Kamper



Fast Fourier transform (FFT)

N-point DFT:

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j2\pi kn/N}$$

$$\left\{ \begin{array}{l} X[0] = \dots \\ X[1] = \dots \\ X[2] = \dots \\ X[3] = x[0] \cdot e^{-j2\pi 3 \cdot 0/N} \\ \quad + x[1] \cdot e^{-j2\pi 3 \cdot 1/N} \\ \quad + \dots \\ \quad + x[N-1] \cdot e^{-j2\pi 3 \cdot (N-1)/N} \\ X[4] = \dots \\ \vdots \\ X[N-1] = \dots \end{array} \right.$$

Complex multiplications:

- for each k in $X[k]$: N complex mults
- for all N values of $X[k]$: N^2 complex mults

radix-2 FFT : Require N to be power of 2

The FFT and DFT lead to precisely the same result. They differ only in the way that they are calculated.

Decimation in time

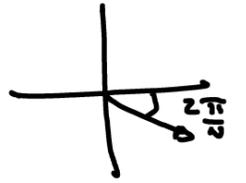
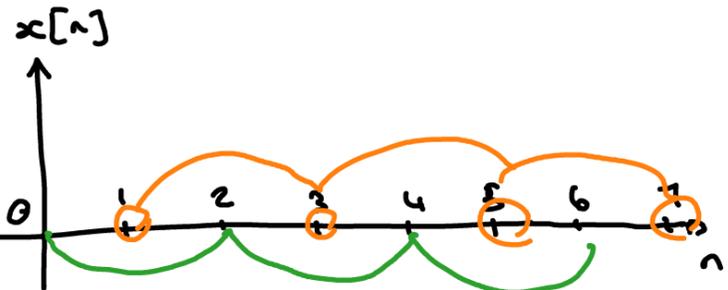
$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j2\pi kn/N}$$

$$= \underbrace{\sum_{n=0}^{N/2-1} x[2n] \cdot e^{-j2\pi k(2n)/N}}_{\text{green bracket}} + \underbrace{\sum_{n=0}^{N/2-1} x[2n+1] \cdot e^{-j2\pi k(2n+1)/N}}_{\text{orange bracket}}$$

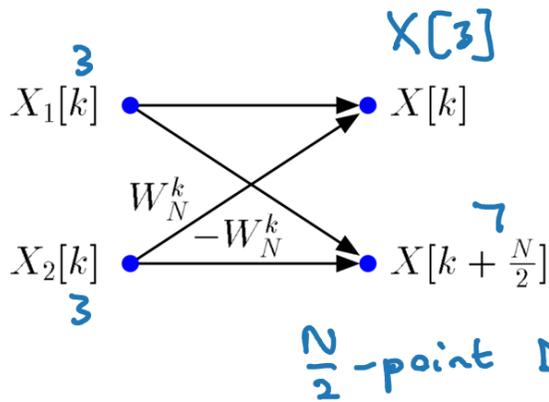
$$= \sum_{n=0}^{N/2-1} x[2n] \cdot e^{-j2\pi k n / N/2} + e^{j2\pi k / N/2} \cdot \sum_{n=0}^{N/2-1} x[2n+1] \cdot e^{-j2\pi k n / N/2}$$

$$= X_1[k] + e^{-j2\pi k / N/2} \cdot X_2[k]$$

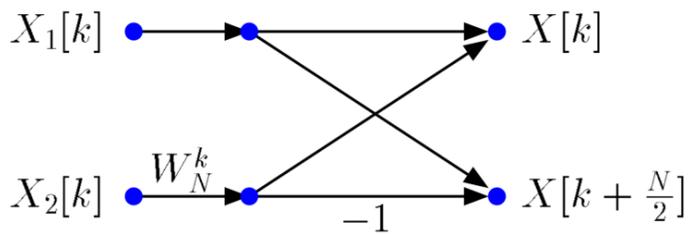
$$= X_1[k] + W_N^k \cdot X_2[k]$$



$$W_N = e^{-j2\pi / N}$$

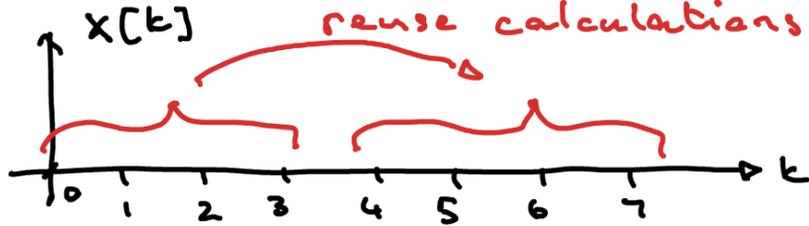


$\frac{N}{2}$ -point DFTs



Complex mults: $2 \times \left(\frac{N}{2}\right)^2 = \frac{N^2}{2}$

Split $X[k]$ calculation in two:



$$X[k] = X_1[k] + W_N^k X_2[k] \text{ for } k=0, 1, \dots, \frac{N}{2}-1$$

$$X[k + \frac{N}{2}] = X_1[k + \frac{N}{2}] + W_N^{k + \frac{N}{2}} X_2[k + \frac{N}{2}] \text{ for } k=0, 1, \dots, \frac{N}{2}-1$$

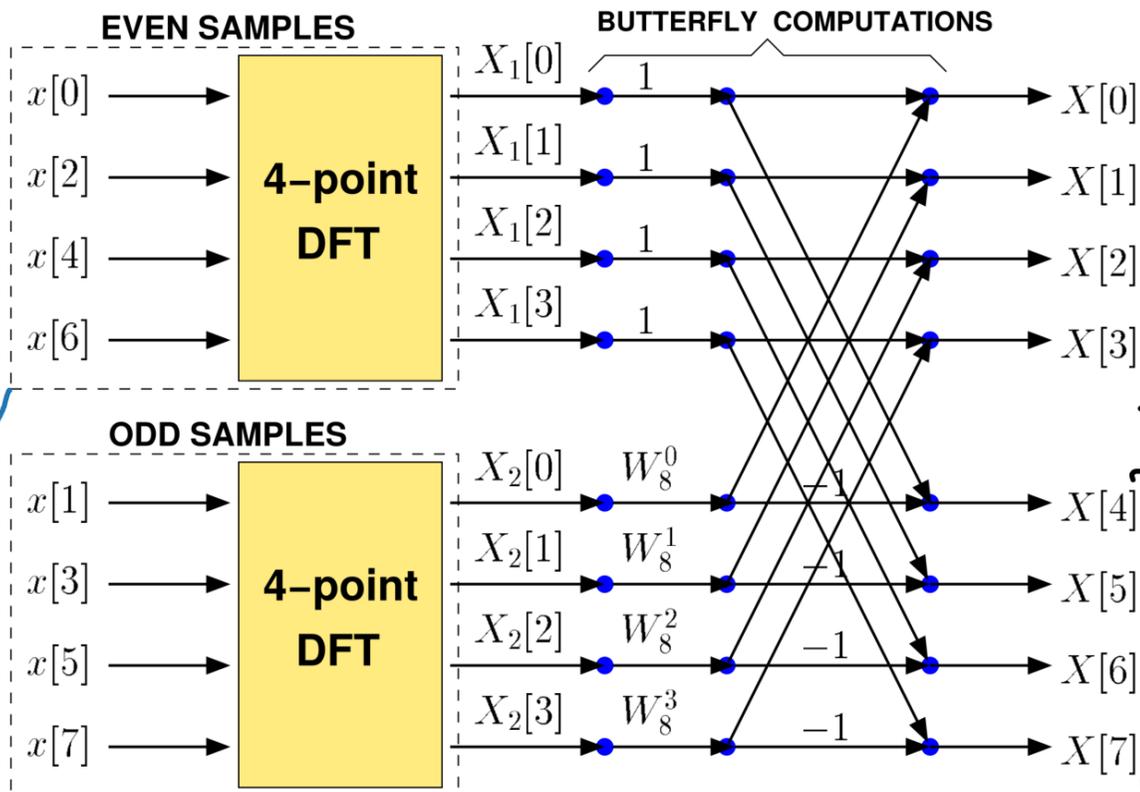
$$= X_1[k] + W_N^{k + \frac{N}{2}} X_2[k]$$

$(e^{-j\frac{2\pi}{N} \cdot \frac{N}{2}}) = e^{-j\pi} = -1$

$$= X_1[k] + W_N^k \cdot W_N^{\frac{N}{2}} X_2[k]$$

$$= X_1[k] - \underbrace{W_N^k}_{\left(\frac{N}{2}\right)^2} X_2[k]$$

Eight-point DFT

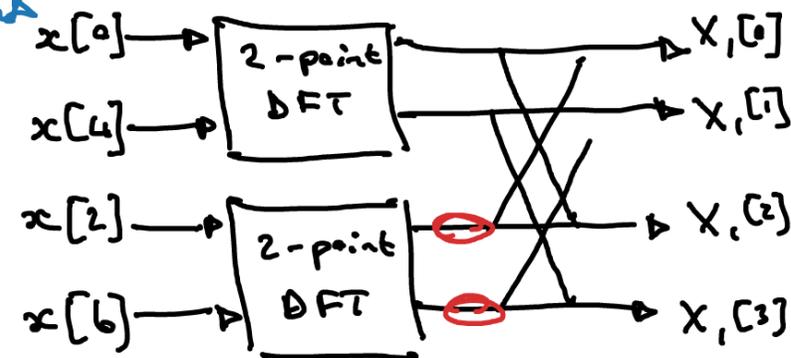


$$X[k] = X_1[k] + W_2^k X_2[k]$$

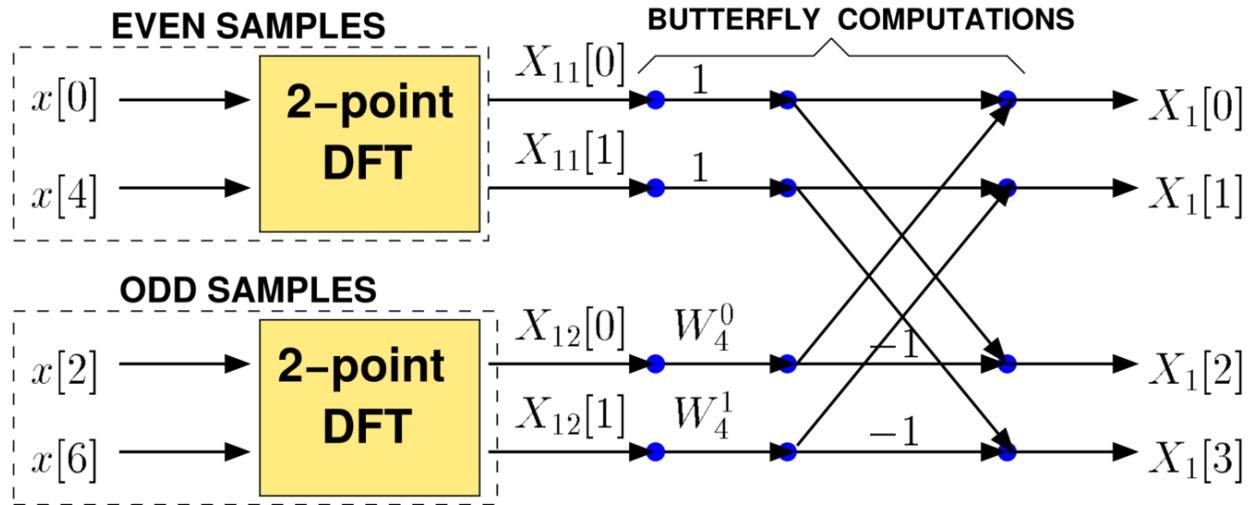
for $k = 0, 1, \dots, \frac{N}{2} - 1$

$$X[k + \frac{N}{2}] = X_1[k] - W_2^k X_2[k]$$

for $k = 0, 1, \dots, \frac{N}{2} - 1$



Four-point DFT

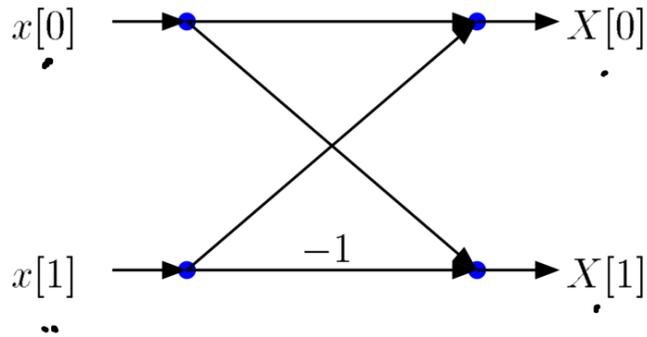


Two-point DFT

$$X[k] = \sum_{n=0}^1 x[n] \cdot e^{-j2\pi kn/2}$$
$$= x[0] \cdot e^0 + x[1] \cdot e^{-j\pi k}$$

$$X[0] = x[0] + x[1]$$

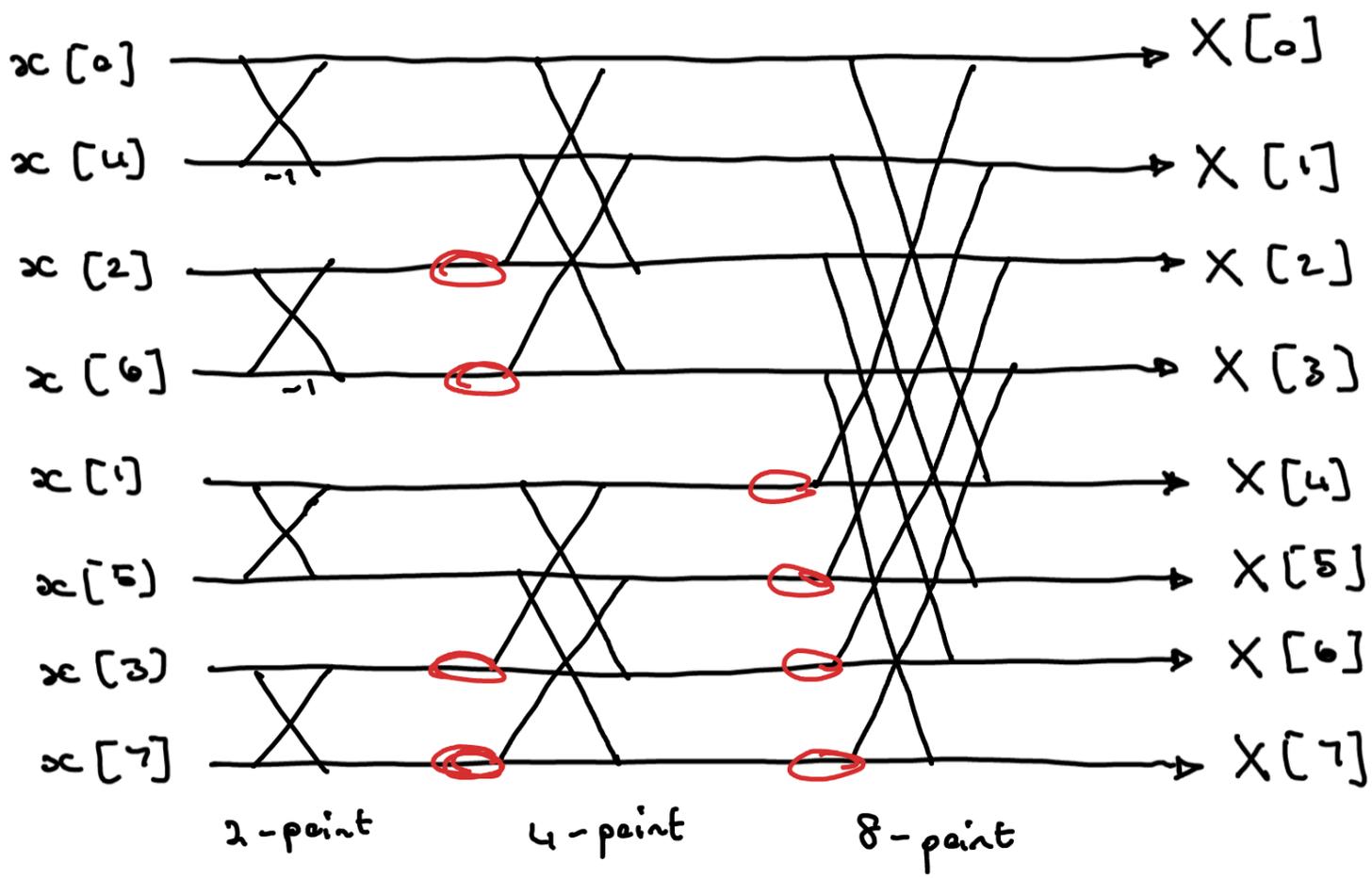
$$X[1] = x[0] - x[1]$$



$$2^{\text{layers}} = N$$

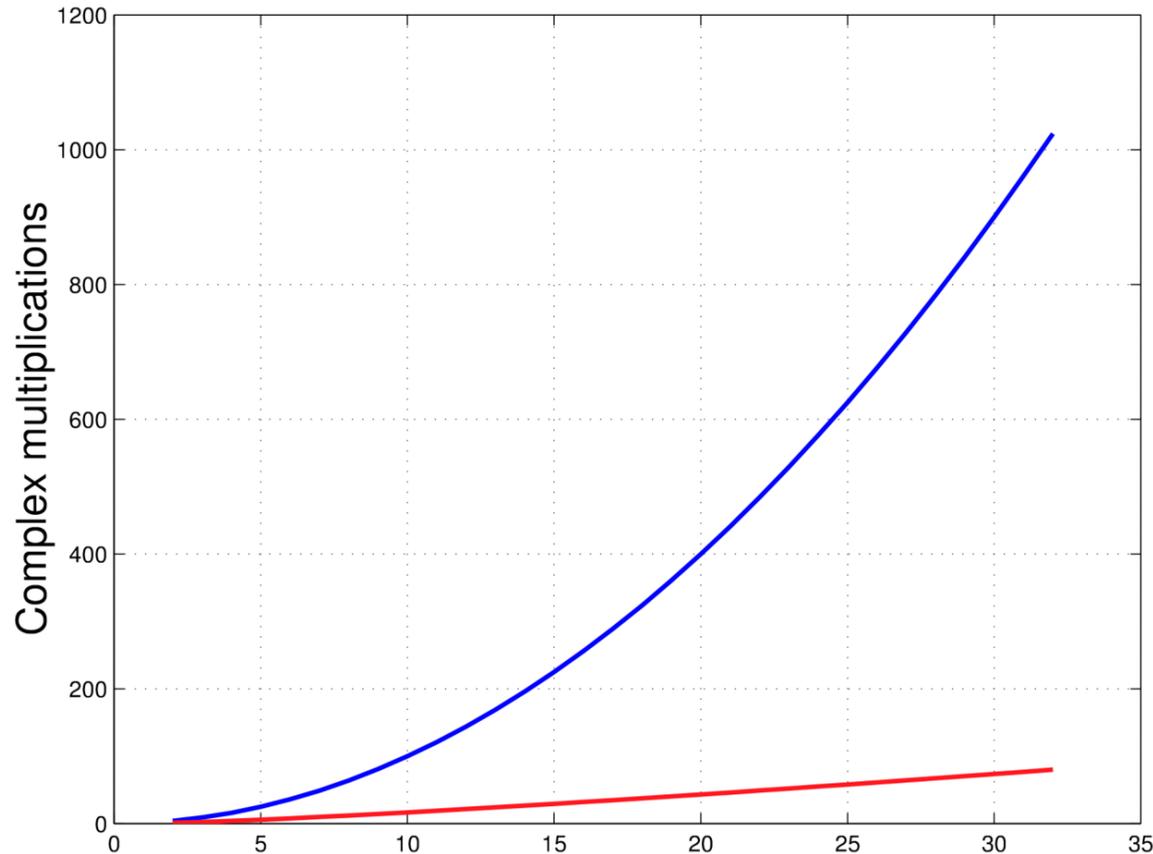
$$\text{layers} = \log_2 N$$

$\frac{N}{2}$ complex mults per layer



$x[2,7]$

Computational complexity



N -point DFT:

- $\log_2 N$ layers

- $\frac{N}{2}$ complex mult per layer

$\therefore \frac{N}{2} \log_2 N$ complex mults
for FFT

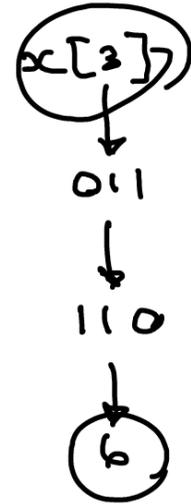
vs

N^2 for direct DFT

1024-point DFT: 1024 ^N 512 complex mults vs. 5120 for FFT

In-place computation of the FFT

Decimal index	Binary index	Bit-reversed index
0: $x[0]$	$x[000]$	000
1: $x[4]$	$x[100]$	001
2: $x[2]$	$x[010]$	010
3: $x[6]$	$x[110]$	011
4: $x[1]$	$x[001]$	100
5: $x[5]$	$x[101]$	101
6: $x[3]$	$x[011]$	110
$x[7]$	$x[111]$	111



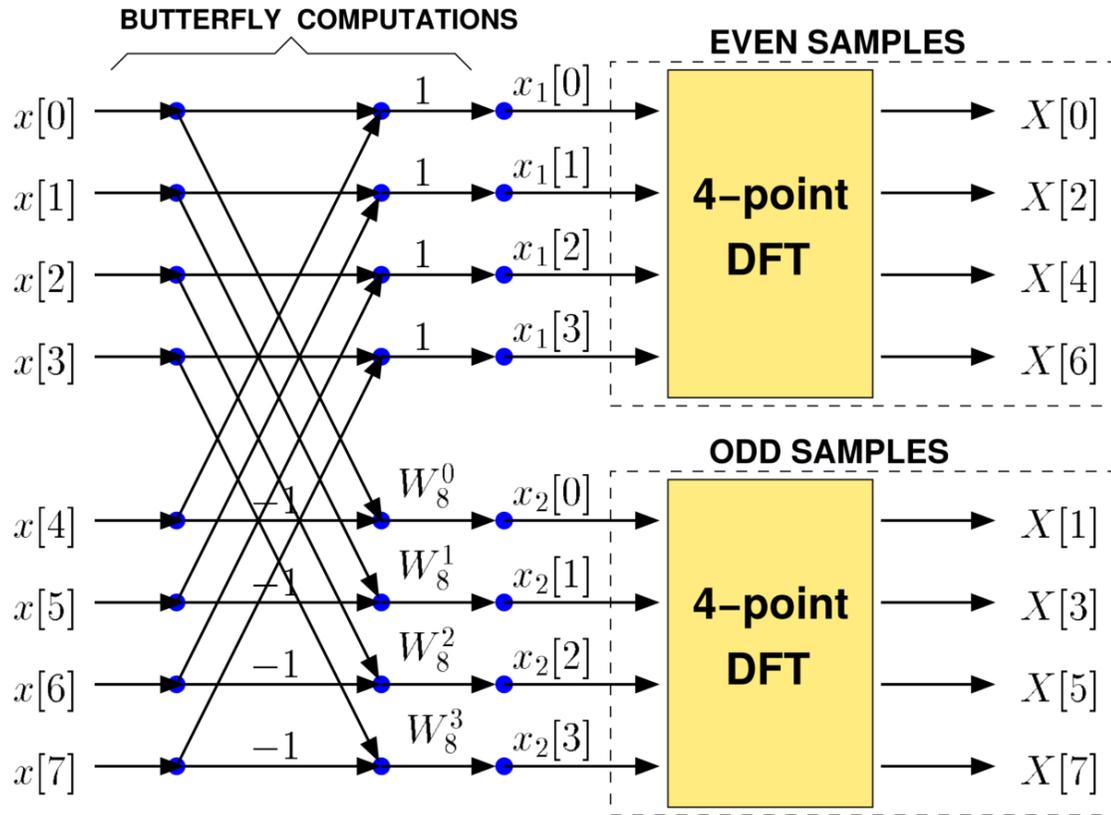
$$x[217] = ?$$

Decimation in frequency

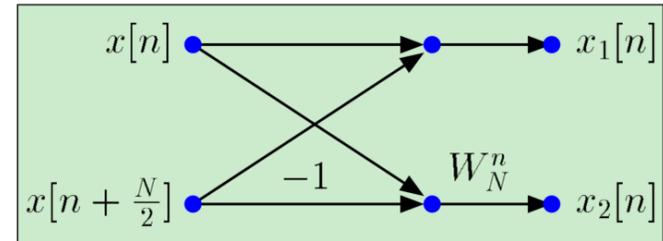
Split $X[k]$ into even and odd sequences:

$$X[2k] \text{ for } k=0, 1, \dots, \frac{N}{2}-1$$

$$X[2k+1] \text{ for } k=0, 1, \dots, \frac{N}{2}-1$$



DECIMATION-IN-FREQUENCY BUTTERFLY COMPUTATION



Summary

Decimation in time:

- Split $x[n]$ into even and odd parts
- Helpers: $X_1[k]$ and $X_2[k]$
- Input: Bit-reversed ; Output: Nice

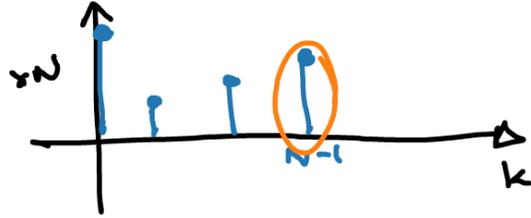
Decimation in frequency:

- Split $X[k]$ into even and odd parts
- Helpers: $x_1[n]$ and $x_2[n]$
- Input: Nice ; Output: Bit-reversed

Calculating IFFTs using FFTs

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j2\pi kn/N}$$

- $\text{FFT} \{X[n]\} = N \cdot x[N-k]$

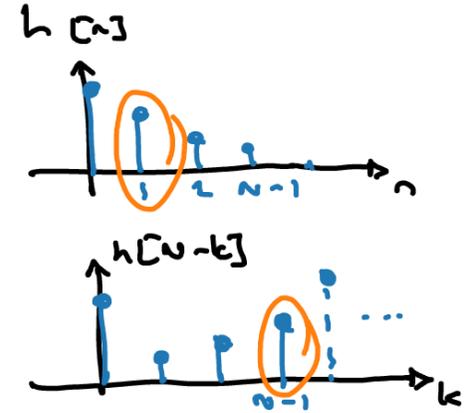


Symmetry for DFT:

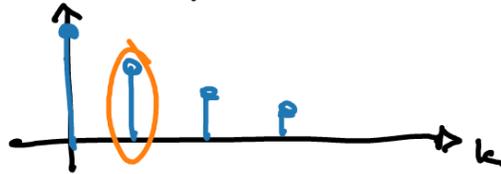
if $\text{DFT} \{h[n]\} = H[k]$

then

$$\text{DFT} \{H[n]\} = N \cdot h[-k] \\ = N \cdot h[\underline{N-k}]$$



- Swap the indices: what happens at $N-k$ now happens at k



- k becomes n
- Divide by N

What to take away

- Some technical tricks
- Divide and conquer: Store intermediate results and use them later
- You need to understand implications of the tricks:
 - When can you (not) use in-place computations?
 - Why radix-2?
 - Why are things slower with `np.fft.fft` when the input is not radix-2?